Easy Geo-Redundancy with MARS + systemd

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LCA2020 Sysadmin Miniconf Presentation by Thomas Schöbel-Theuer

Easy Geo-Redundant Handover+Failover: Agenda



- Short background: MARS for end-to-end SLA 99.98%
- Cluster management for long distances
- Using systemd as a clustermanager
- **Current Status / Future Plans**

Background: SLA + Growth at 1&1 Ionos ShaHoLin



= Shared Hosting Linux

- SLA: 99.98% end-to-end measured from Frankfurt
 - Including WAN outages, PHP problems, HumanError™
 - => MARS geo-redundancy must compensate much better!
- 4 datacenters at 2 continents, pair distance > 50 km
 - ~ 9 millions of customer home directories
- ~ 10 billions of inodes + daily incremental backup
- > 4.7 petabytes *allocated* in ~ 3800 xfs instances
 - LocalStorage LVM ~ 8 PB x 2 for geo-redundancy via MARS
 - https://github.com/schoebel/mars
- Growth rate ~ 21 % / year
- **Solution: Container Football on top of MARS**

https://github.com/schoebel/football

Cluster Management for > 50 km



- Proprietary e.g. 1&1 cm3 (no GPL)
 - Pacemaker & co typically don't work as expected
 - original HeartBeat DSM model: shared disk cannot really handle Split Brain

explainable by CAP theorem

- Using systemd as a Linux clustermanager
 - already in use almost everywhere e.g. startup of VMs
 - itself somewhat "monolithic", but extensible via Unit Files
 - path watchers can monitor /mars/resource-\$res/* remote updates from MARS cluster communication => generic remote control
- MARS dynamic resource creation / deletion marsadm join-resource / leave-resource
- Solution: marsadm internal macro processor creates / deletes systemd units "on the fly"

systemd Unit Example (Template)



```
bash> cat \^\{mntname\}-@\{res\}.mount
@eval{%let{mntpath}{%subst{%{mntname}}{-}{/}}}
[Unit]
Description=MARS local mount on /@{mntpath}/@{res}
Documentation=https://github.com/schoebel/mars/docu/mars-user-manual.pdf
Requires=mars.service
                                               Use MARS' metadata symlink
After=mars.service
                                                updates for remote control
ConditionPathIsSymbolicLink=/mars/resource-@{res}/systemd-want
ConditionPathExists=/mars/resource-@{res}/userspace/systemd-want-@{host}
ConditionPathExists=/dev/mars/@{res}
ConditionPathIsDirectory=/@{mntpath}/@{res}
[Mount]
What=/dev/mars/@{res}
Where=/@{mntpath}/@{res}
[Install]
WantedBy=mars.service
```

Usage of systemd unit templates



- Activation of template (once after resource creation, for the whole cluster)
 marsadm create-resource \$resource /dev/\$vg/\$resource
 mkfs.xfs /dev/mars/\$resource
 marsadm set-systemd-unit \$resource \$start_unit \$stop_unit
 => automatic instantiation via macro processor
- Usage at planned handover: piggyback on distributed MARS symlinks marsadm primary \$resource (or marsadm primary all)
 - Automagically (independently for each resource):
 - Old primary: systemctl stop \$stop_unit
 - Old primary: MARS goes to secondary mode
 - New primary: MARS becomes primary /dev/mars/\$resource will appear
 - New primary: systemctl start \$start_unit
- **Usage at unplanned failover:**
 - marsadm disconnect all ; marsadm primary --force all

Current status of systemd interface



- BETA feature! not yet in production
 - example templates in systemd/ subdir
- **Currently works sequentially**
 - observation: systemctl is non-reentrant, can deadlock
 - marsadm uses (breakable) locks for protection
- Planned improvements: all resources in parallel to each other
 - Needs heavy testing
 - Help from the community welcome!
 - e.g. contribute new systemd templates for KVM startup, or iSCSI / NFS exports, ...

MARS Current Status



MARS source under GPL + docs:

github.com/schoebel/mars
docu/mars-user-manual.pdf

docu/mars-architectureguide.pdf

- mars0.1stable productive since 02/2014
- Backbone of the 1&1 Ionos geo-redundancy feature
 - up to 14 LXC Containers on 1 Hypervisor
 - Efficiency project using Football:
 - TCO has **halved!**



MARS Future Plans

Kernel part almost done mixed operations of old/new MARS versions



Faster checksumming (CRC32 | CRC32 | SHA1 | MD5)

Logfile compression (LZO | LZ4 | ZLIB)

Optional network transport compression

may help for some very slow networks

IO data paths already scaling well

TODO: better *metadata* scalability needed!

- single mars_main control thread (non-blocking)

- TODO: more resources per host (max. 24 in prod at 1&1)

TODO: more hosts per cluster

TODO: Linux kernel upstream

- requires a *lot* of work!
- tomorrow's presentation at kernel miniconf

TODO: more tooling, more systemd templates, integration into other OpenSource projects, ...

Collaboration sought

=> Opportunities for other OpenSource projects!



Sponsoring (MARS + Football)

- Best for > 1 PiB of enterprise-critical data
 - More Football plugins in future, e.g. for KVM, ...
- Future pool-optimizer will deliver similar functionality than **Kubernetes**
 - but on stateful storage + containers instead of stateless Docker containers
 - State is in the storage and in the machines, but not in orchestration
- Long-term perspective
 - MARS is largely complementary to DRBD
 - Geo-redundancy with OpenSource components
 - distances > 50km possible, tolerates flaky replication networks
 - Price / performance better than anything else (see mars-achitecture-guide.pdf)
 - Architectural reliability better than BigCluster with cheaper hw + network!
- ask me: decades of experience with enterprise-critical data and long-distance replication



Why GEO-Redundancy

DR = Disaster Recovery CDP = Continuous Data Protection



- Example: GALILEO incident (DR / CDP did not work)
 - Disaster = earthquake, flood, terrorist attack, power outage, ...
- **BSI Paper 12/2018:**

Kriterien für die Standortwahl höchstverfügbarer und georedundanter Rechenzentren

https://www.bsi.bund.de/SharedDocs/Downloads/DE/BSI/Sicherheitsberatung/Standort-Kriterien_HV-RZ/Standort-Kriterien_HV-RZ.pdf?__blob=publicationFile&v=5

in English: Criteria for Locations of Highly Available and Geo-Redundant Datacenters

- Stimulated some controversial discussions, but see commentary https://www.it-finanzmagazin.de/bsi-rechenzentren-entfernung-bafin-84078/
- Conclusions: distances > 200 km "recommended"
 - Might influence future legislation (EU / international)
 - "Critical Infrastructures" more important!

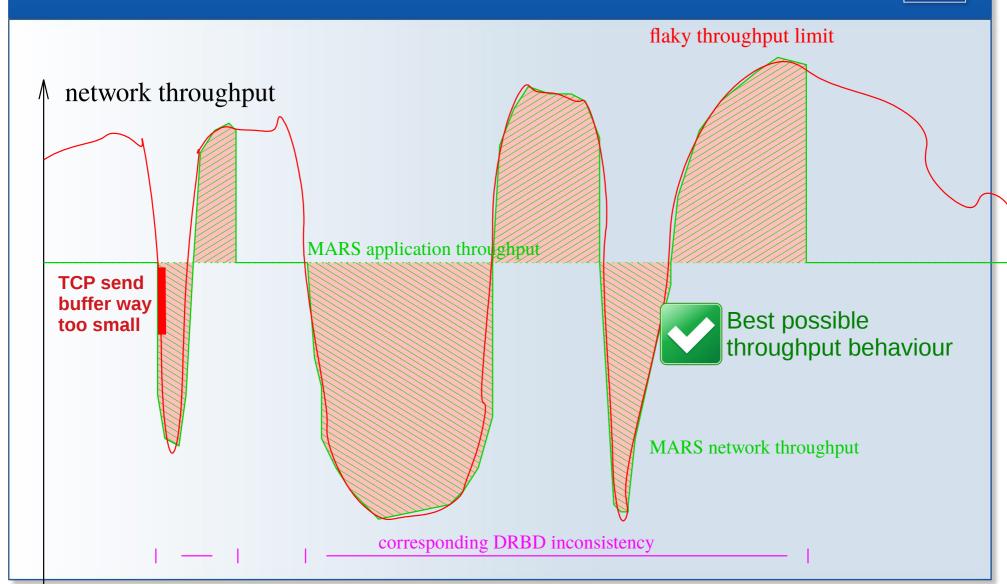
Long-Distance Asynchronous Replication



- Synchronous does not *generally* work over ≈50 km
 - like iSCSI over 50 km
- Need Asynchronous Replication
 - Application specific, e.g. mySQL replication
 - Commercial appliances: \$\$\$ €€€
 - OpenSource
 - plain DRBD is NOT asynchronous
 - commercial DRBD-Proxy: RAM buffering
 - MARS: truly asynchronous + persistent buffering
 - + transaction logging + MD5 checksums
 - + Anytime Consistency

Network Bottlenecks: MARS





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time

Replication at Block Level vs FS Level





Apache, PHP, mySQL, Mail Queues, etc

Potential Cut Point A for Distributed System

~ 25 Operation Types

~ 100.000 Ops / s



--- NO long distances

xfs, ext4, btrfs, zfs, ... vs nfs, Ceph, Swift, ...



Potential Cut Point B for Distributed System

DSM = Distributed Shared Memory

=> Cache Coherence Problem!

Caching Layer

Kernelspace

1:100 reduction

dentry Cache, ...

Page Cache,

2 Operation Types (r/w) ~ 1.000 Ops / s

Block Layer

+++ LONG DISTANCES

Hardware-RAID, BBU, ...

LVM, DRBD / MARS >

Potential Cut Point C for Distributed System

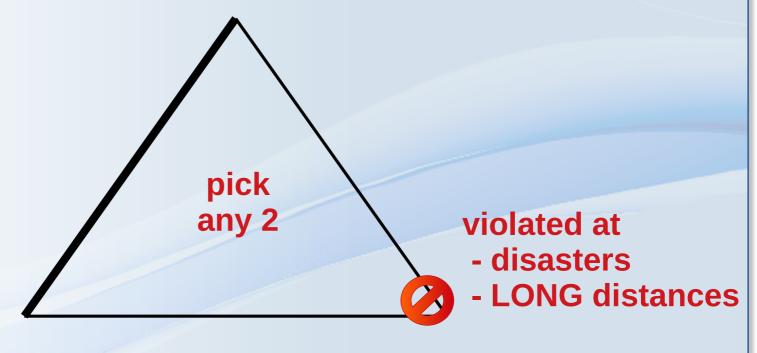
+++ replication of VMs for free!

Hardware

CAP Theorem







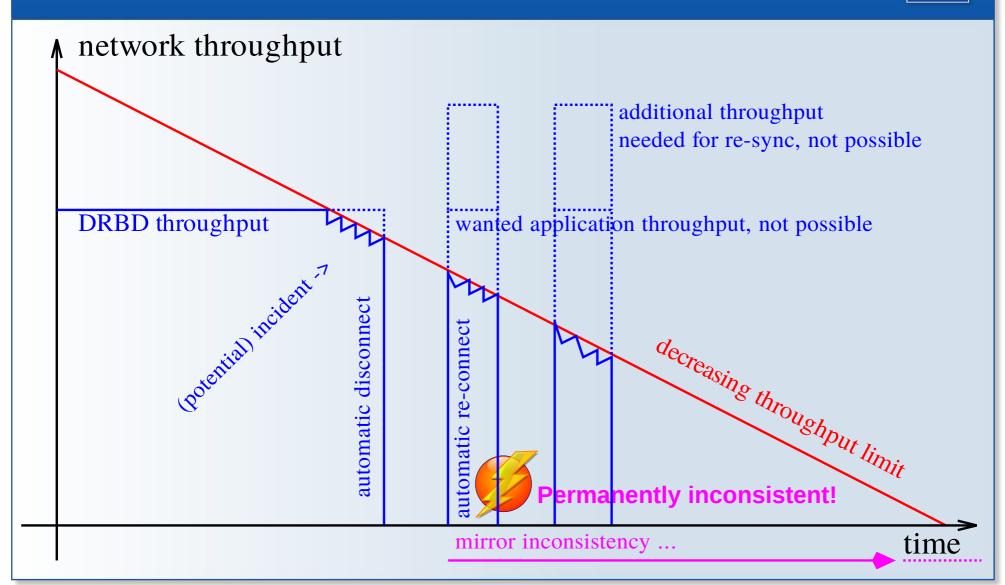
A = Availability

P = Partitioning Tolerance

= the network can have its own outages

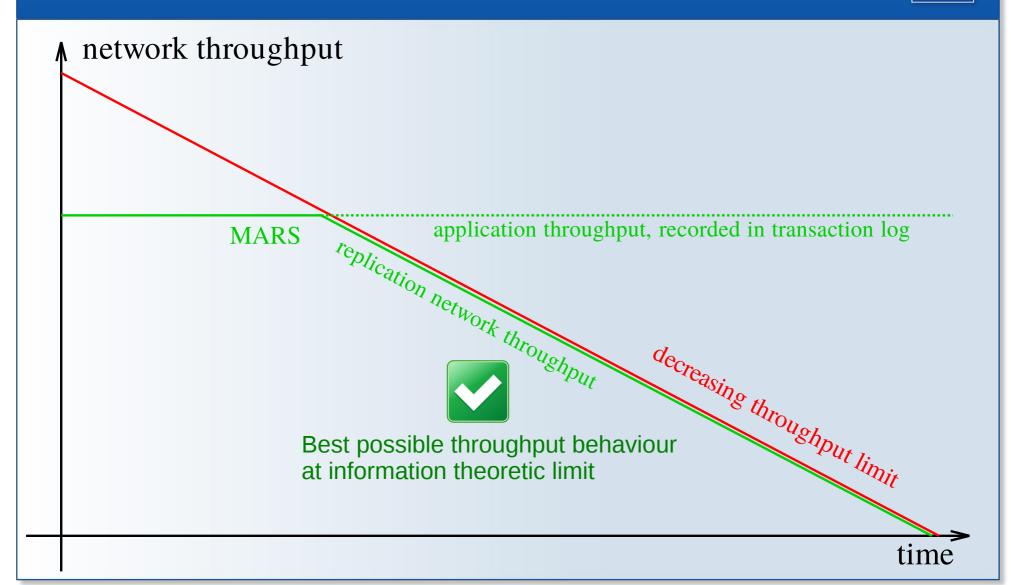
Network Bottlenecks (1) DRBD





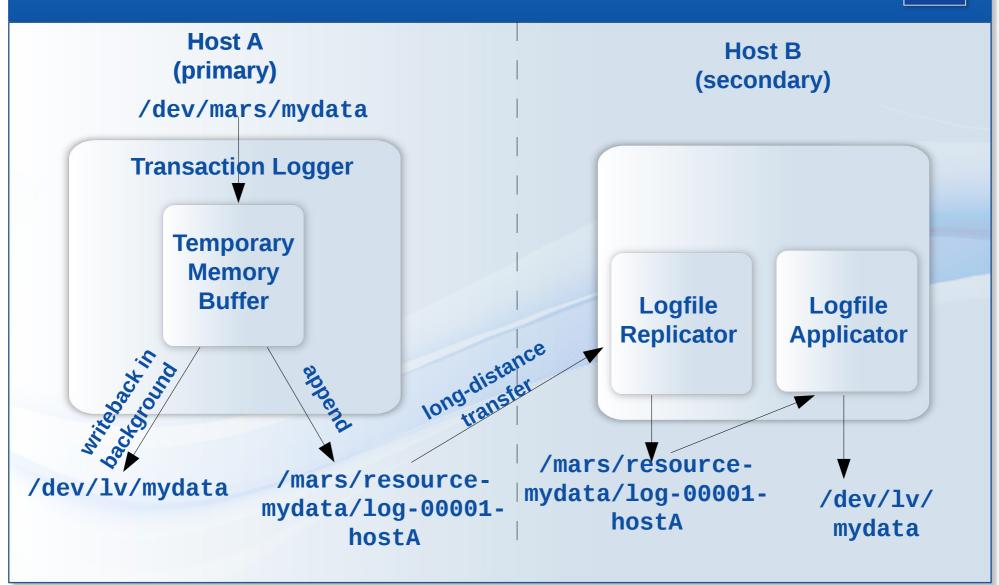
Network Bottlenecks (2) MARS





MARS Data Flow Principle





Use Cases DRBD+proxy vs MARS



DRBD+proxy (proprietary)

Application area:

- Distances: any
- Aynchronously
 - Buffering in RAM
- Unreliable network leads to **frequent re-syncs**
 - RAM buffer gets lost
 - at cost of actuality
- **Long** inconsistencies during re-sync
- Under pressure: **permanent** inconsistency possible
- High memory overhead
- Difficult scaling to k>2 nodes

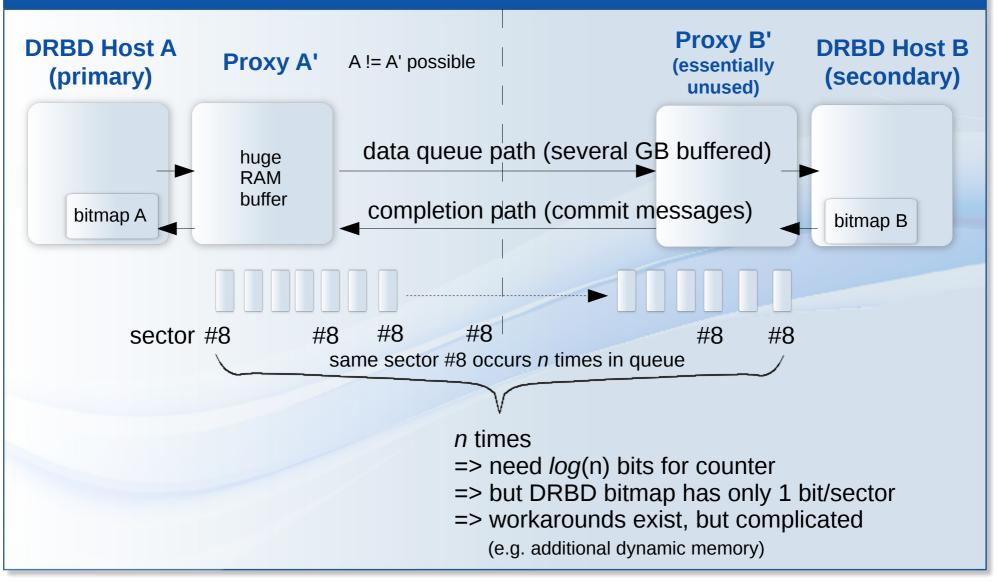
MARS (GPL)

Application area:

- Distances: any (>>50 km)
- Asynchronously
 - near-synchronous modes in preparation
- Tolerates unreliable network
- Anytime consistency
 - no re-sync
- Under pressure: no inconsistency
 - possibly at cost of actuality
- Needs >= 100GB in /mars/ for transaction logfiles
 - dedicated spindle(s) recommended
 - RAID with BBU recommended
- Easy scaling to k>2 nodes

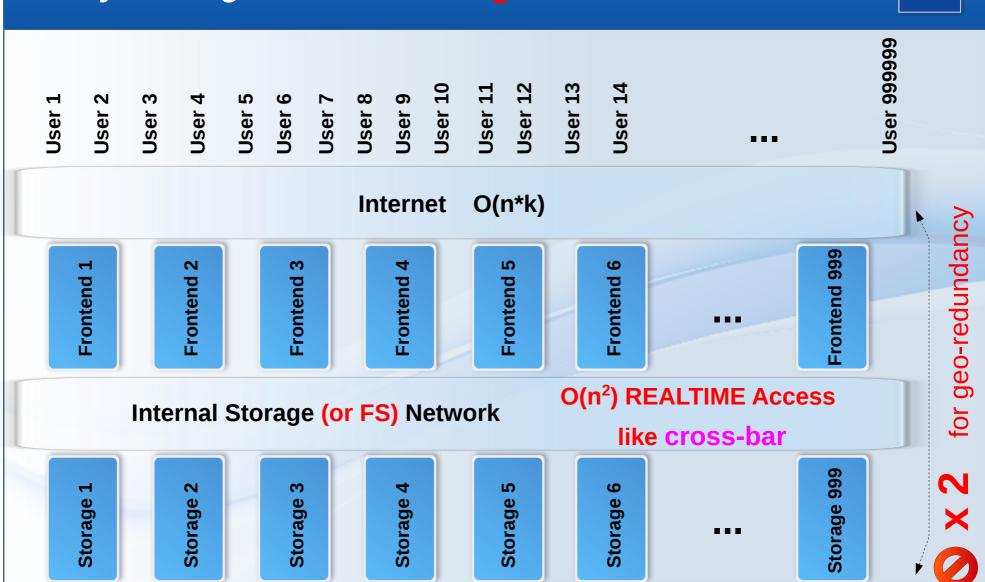
DRBD+proxy Architectural Challenge





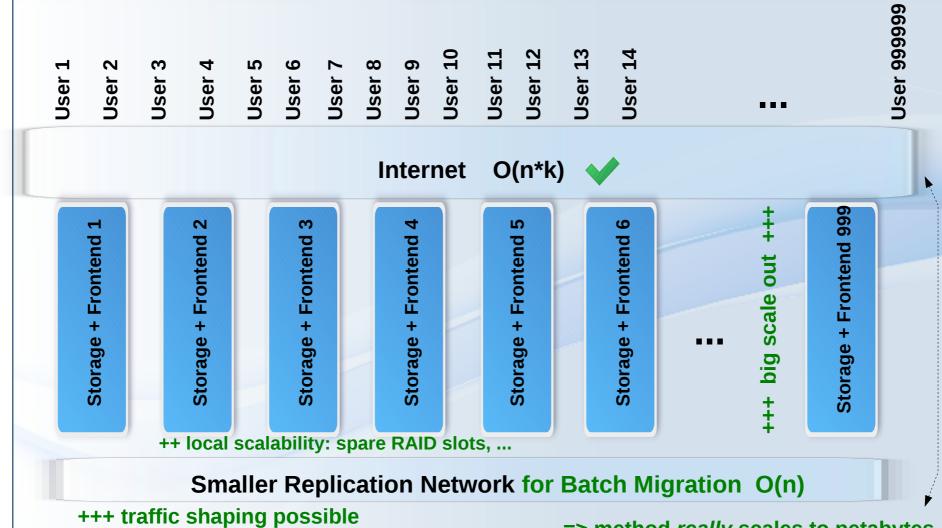
Badly Scaling Architecture: Big Cluster





Well-Scaling Architecture: Sharding





=> method *really* scales to petabytes

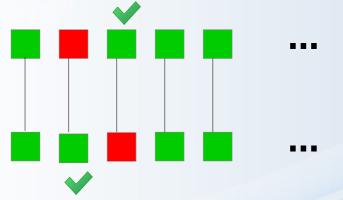
geo-redundancy

Reliability of Architectures: NODE failures



2 Node failure => ALL their disks are unreachable

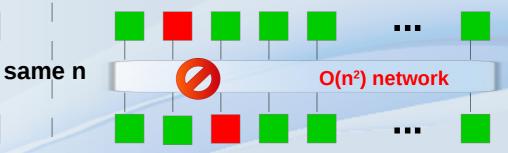




=> no customer-visible incident

Low probability for hitting the *same* pair, even then: only 1 shard affected => low total downtime

Big Storage Cluster e.g. Ceph, Swift, ...



k=2 replicas not enough
=> INCIDENT because objects are randomly
distributed across whole cluster

Higher probability for hitting any 2 nodes, then O(n) clients affected => much higher total downtime

need k >= 3 replicas here

Cost (1) non-georedundant, n>100 nodes



- Big Cluster:
 Typically ≈RAID-10 with k=3 replicas for failure compensation
- **Disks:** > 300%
- Additional CPU and RAM for storage nodes
- Additional power
- Additional HU

- Simple Sharding:
 Often local RAID-6
 sufficient (plus external backup, no further redundancy)
- **Disks: < 120%**
- Client == Server
 no storage network

MARS for LV background migration

- Hardware RAID controllers
 with BBU cache on 1 card
- Less power, less HU

Cost (2) geo-redundant => LONG Distances



- Big Cluster:
 - 2X ≈ RAID-10 for failure compensation
 (k=6 replicas needed, smaller does not work in long-lasting DC failure scenarios)
- **Disks: > 600%**
- Additional CPU and RAM for storage nodes
- Additional power
- Additional HU

- **Geo-redundant Sharding:**
 - 2 x local RAID-6
 - MARS for long distances
 or DRBD for room redundancy
- **Disks:** < 240%
- Hardware RAID controllers with BBU
- Less power
- Less HU

Cost (1+2): Geo-Redundancy Cheaper than Big Cluster

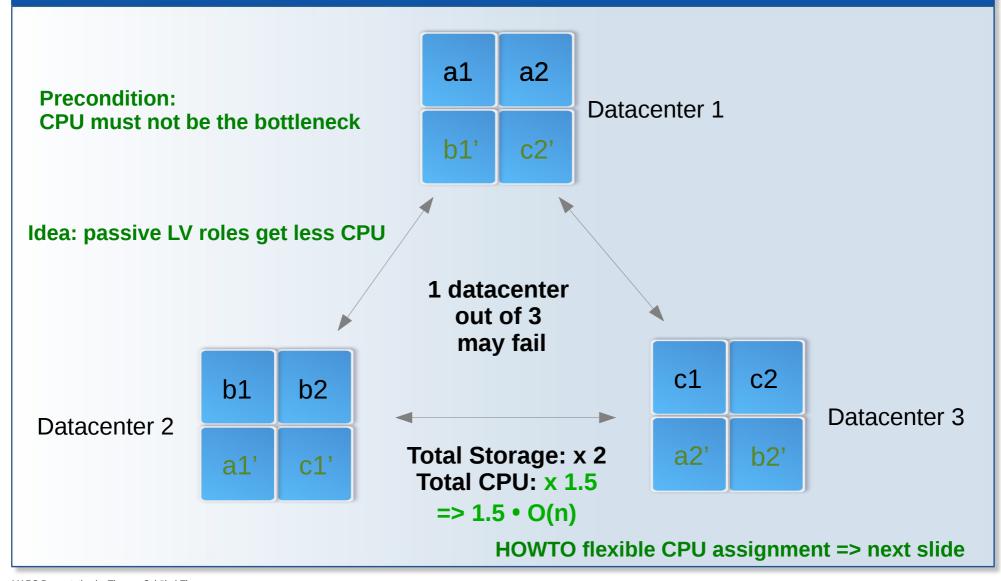


- **Single Big Cluster:**
 - − ≈RAID-10 with k=3 replicas for failure compensation
- O(n) Clients
 - + 3 O(n) storage servers
 - + O(n²) storage network
- **Disks: > 300%**
- Additional power
- Additional HU

- **Geo-redundant sharding:**
 - 2 x local RAID-6
 - MARS for long distances
 or DRBD for room redundancy
- 2 O(n) clients = storage servers
 - + O(n) replication network
- **Disks:** < 240%
- Less total power
- Less total HU
 +++ geo failure scenarios

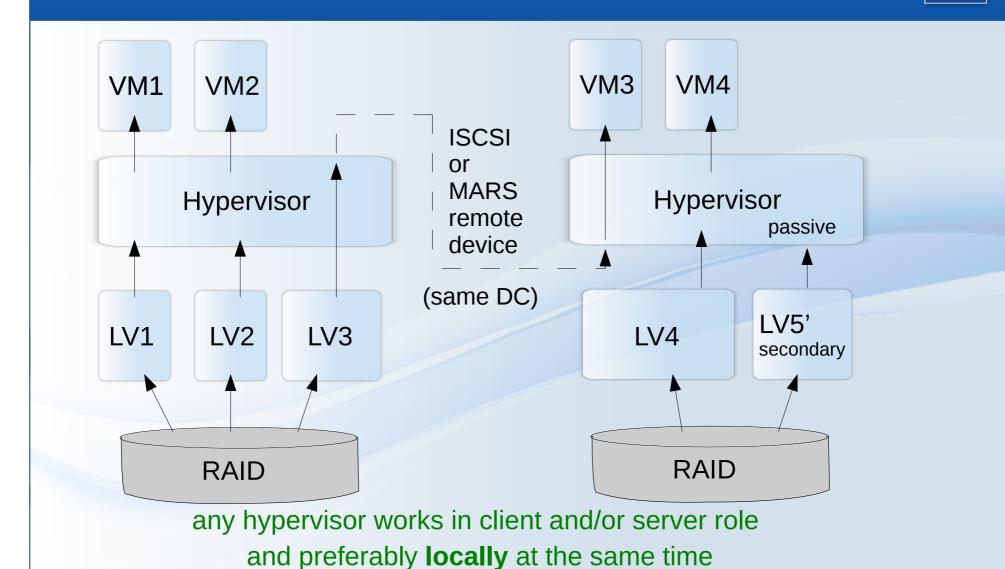
Costs (3): Geo-Redundancy even Cheaper





Flexible MARS Sharding + Cluster-on-Demand

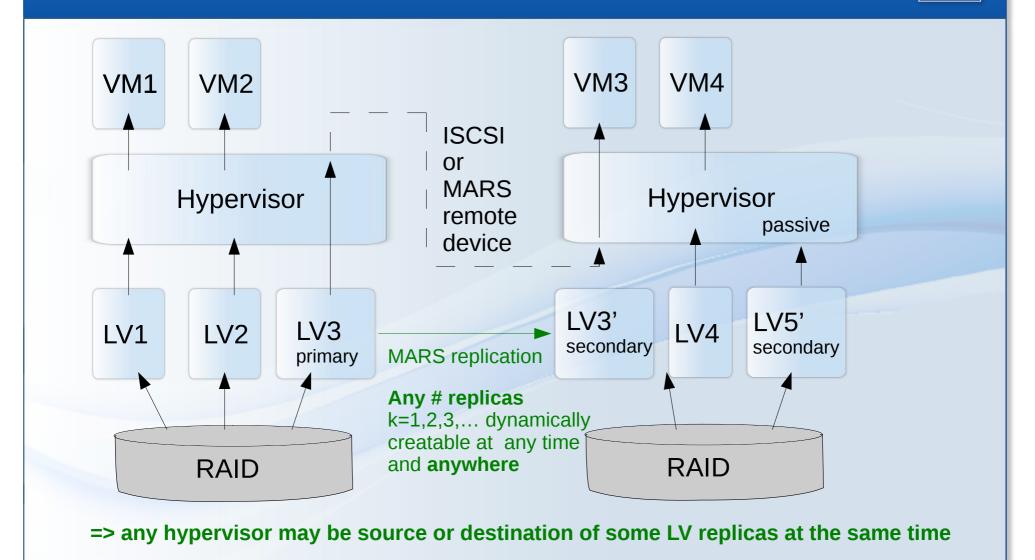




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Flexible MARS Background Data Migration





Football Current Status



- GPL with lots of plugins, some generic, some 1&1-specific
 - about 2/3 of code is generic
 - plugins/football-basic.sh uses systemd as cluster manager
 - https://github.com/schoebel/football
 - https://github.com/schoebel/mars
- Multiple operations:
 - migrate \$vm \$target_cluster
 - low downtime (seconds to few minutes)
 - shrink \$vm \$target_percent
 - uses local incremental rsync, more downtime
 - expand \$vm \$target_percent
 - online, no downtime
- In production at 1&1 lonos
 - get rid of old hardware (project successfully finished)
 - load balancing
 - >50 "kicks" per week
 - limited by hardware deployment speed
 - Proprietary Planner (for HW lifecycle)

